# ON THE OPTIMAL CHOICE OF MEASUREMENTS IN LINEAR QUADRATIC GAUSSIAN TEAM PROBLEMS

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### Abstract

The problem of optimal choice of information for some simple linear quadratic gaussian team problems is considered. The unknowns to be chosen subject to constraints are the matrices involved in the linear measurements available to the decision makers. For several types of such problems, characterizations of the best choices of these matrices are given and several results illustrating the meaning of these characterizations and ways for finding the optimal choices are also presented.

#### · 1. Introduction

The purpose of this paper is to examine the problem of optimal choice of information in a simple stochastic team set-up. Most of the papers which consider stochastic optimization problems, assume that the information used is provided by given measurement devices and try to characterize or find the optimal control laws. Nonetheless, the information to be used--or the measuring devices--is quite often subject to the choice of the decision makers, and although the decision makers would wish to have all the possible information available, it could oftentimes be an impossible burden to collect and (or) process it. Thus, a decision maker is obliged to choose, subject to constraints, what information he will use. The choice of information might be a nontrivial problem, because the objective to be achieved does not always reveal in a straightforward manner what the essentially needed information is. Having such considerations in mind we formulate some problems related to the choice of information. Our objective functions are quadratic, the measurements linear and the random variables involved gaussian. We try to characterize the best choice of information subject to restrictions which usually assume the form of upper bounds on the rank of some matrices and can be interpreted as restrictions on the number of linearly independent measurements available to the decision makers. We also impose occasionally the condition that the information of two different decision makers are orthogonal.

Problems related to the optimal choice of information have been previously considered in several papers as for example in [1, 2, 3, 4, 5]. In [1] the problem of finding the best measurement in order to achieve the minimum possible value of a quadratic cost is considered for a static problem. In contrast with [2, 3, 4] where the cost is the covariance of the error of the Kalman estimate, [1] considers an arbitrary quadratic cost. Our framework is very similar to the one of [1], but our attention is directed to the case where there are more than one decision makers, i.e., we deal with a team problem. In [1], an algorithm for finding the optimal choice of information for the case of two or more decision makers is suggested but as is pointed out in [1] and also demonstrated by an example in Section 3, this algorithm might fail to converge to the global solution; this is essentially due to the nonconvex character of the underlying optimization problem. In the present paper we characterize the optimal choice of information in terms of "generalized type eigenvalue" problems,

which at present seem quite nontrivial to solve in their generality. Some examples are also considered in order to elaborate on the difficulties associated with solving such problems. Some of the work presented have appeared in a preliminary form in [5].

#### 2. Problem Statement

Let x be a Gaussian random vector in  $\mathbb{R}^n$  with zero mean and unit variance. The measurements  $y_1$ ,  $y_2$  are defined by

$$y_1 = C_1 x$$

$$y_2 = C_2 x$$
(1)

where  $C_1$ ,  $C_2$  are real constant matrices of dimensions  $r_1 \times n$ ,  $r_2 \times n$  respectively. Let  $\gamma_1$ :  $R^{r_1} \rightarrow R^{m_1}$ ,  $\gamma_2$ :  $R^{r_2} \rightarrow R^{m_2}$  be two functions and set

$$u_1 = Y_1(y_1)$$
 $u_2 = Y_2(y_2)$  (2)

 $\gamma_1, \gamma_2$  are chosen as to minimize the cost

$$J(\gamma_{1}, \gamma_{2}) = E\left[\frac{1}{2}u'_{1}u_{1} + \frac{1}{2}u'_{2}u_{2} + u'_{1}Ru_{2} + u'_{1}S_{1}x + u'_{2}S_{2}x\right]$$
(3)

The matrices R,  $S_1$ ,  $S_2$  are real, constant, with appropriate dimensions and it holds

$$\begin{bmatrix} I & R \\ R' & I \end{bmatrix} > 0 \tag{4}$$

i.e., J is strictly convex in  $(u_1, u_2)$ . Of course,  $\gamma_1, \gamma_2$  have to satisfy the appropriate measurability assumptions. It is known that the pair  $(\gamma_1^*, \gamma_2^*)$  which minimizes (3) exists and is of the form

$$\gamma_1^*(y_1) = L_1 y_1, \quad \gamma_2^*(y_2) = L_2 y_2$$
 (5)

where L<sub>1</sub>, L<sub>2</sub> are two matrices satisfying the system of equations

$$\mathbf{L_{1}C_{1}} + \mathbf{R}\,\mathbf{L_{2}C_{2}C_{1}'(C_{1}C_{1}')}^{+}C_{1} + \mathbf{S_{1}C_{1}'(C_{1}C_{1}')}^{+}C_{1} = 0 \quad (6)$$

$$L_{2}C_{2}+R'L_{1}C_{1}C_{2}(C_{2}C_{2})^{+}C_{2}+S_{2}C_{2}(C_{2}C_{2})^{+}C_{2}=0$$
 (7)

(+ denotes the pseudoinverse). Equations (6) and (7) can be solved uniquely for  $L_1C_1$ ,  $L_2C_2$ . The optimal cost  $J^*$  is uniquely determined and if we consider that  $R, S_1, S_2$  are fixed,  $J^*$  can be considered as a function of  $C_1$  and  $C_2$ , i.e.,:

$$J^* = \hat{J}(C_1, C_2)$$
 (8)

Thus one is motivated to consider problems of the

form

minimize 
$$\hat{J}(C_1, C_2)$$

i. e., to consider what is the best choice of information in the sense that the resulting optimal cost is as small as possible.

## 3. Case 1: One Decision Maker

The results presented here overlap to some extent with those in [1], but we include them for reasons of completeness. In the case of one decision maker, the cost and the measurements are given by

$$J = E\left[\frac{1}{2}u^{\dagger}u + u^{\dagger}Sx\right] \quad y = Cx \quad u = \gamma(y) \quad (10)$$

The optimal u is given by

$$u = -SPx = Ly \tag{11}$$

where P is the projection matrix which projects on the range of C', i.e.,

$$P = C'(CC')^{\dagger}C'$$
,  $L = -SC'(CC')^{\dagger}$  (12)

The optimal cost is given by

$$J^* = \hat{J}(C) = E[-\frac{1}{2}x'PS'SPx] = -\frac{1}{2}tr[PS'SP]$$

$$= -\frac{1}{2}tr[S'SP]$$
 (13)

Let us first consider the problem of minimizing  $\hat{J}(C)$  subject to the restriction that no more than  $\rho$  linearly independent measurements should be available, or equivalently rank  $C \leq \rho$ , or equivalently rank  $P \leq \rho$ . Formally:

max tr [S'SP]  
subject to: rank 
$$P \le \rho$$
 (14)

where  $\rho \leq n$ . The solution of Problem (15) is known and can be obtained by taking P to be the projection matrix which projects on the space spanned by the  $\rho$  eigenvectors which correspond to the  $\rho$  largest eigenvalues of S'S.

A slight generalization of the problem considered above is the following: C can be chosen as any matrix with maximal rank p but it has to be of the form

$$C = TC_0 (15)$$

for some matrix T, where  $C_0$  is a given matrix with rank greater or equal than  $\rho$ ; i. e., we essentially have to choose the best  $\rho$  dimensional subspace lying within the space range  $C_0'$ . It is easy to see that if  $P_0$  is the projection matrix on the range of  $C_0$ , this problem can be equivalently stated as

$$\max \operatorname{tr} \left[ P_0 S' S P_0 \cdot P \right] \tag{16}$$

subject to rank  $P \leq \rho$ 

and can be solved in a similar fashion.

In the formulations given above, we assumed that at most  $\rho$  linearly independent measurements of x can be obtained, each one of them with perfect accuracy. If we want to consider the case where we can make at most  $\rho$  linearly independent measurements, but there is a fixed measurement noise, we need a different formulation. Namely, let

$$\mathbf{x} = \begin{bmatrix} \mathbf{x}_1 \\ --- \\ \mathbf{v} \end{bmatrix}, \quad \mathbf{S} = \begin{bmatrix} \mathbf{S}_1 & \mathbf{0} \end{bmatrix}, \quad \mathbf{S}\mathbf{x} = \mathbf{S}_1\mathbf{x}_1 \tag{17}$$

$$C = [C_1 \mid I]$$
  $y = C_1 x_1 + v$  (18)

Only  $x_1$  appears in the cost and v represents the noise in the measurements. For fixed  $C_1$ , the optimal cost is given by (see also (15))

$$\hat{J}(C) = -\frac{1}{2} \operatorname{tr} \left[ S_1' S_1 C_1' (I + C_1 C_1')^{-1} C_1 \right]$$
 (19)

We are going to impose two constraints on  $C_1$ . The first one is rank  $C_1 \leq \rho$ . The second one is  $\|C_1\| \leq a$ , i. e., a magnitude constraint on  $C_1$ , where a is some fixed positive constant. Such a magnitude constraint would represent no restriction to the previously considered cases, since what matters there is only the range of C'. Thus we want to solve

max tr 
$$[S_1'S_1 C_1' (I + C_1C_1')^{-1}C_1]$$
  
subject to: rank  $C_1 \le \rho$  (20)  
 $||C_1|| \le a$ 

(We employ the equal Euclidean norm for vectors and the sup norm for matrices.) To solve (20) we proceed as follows. Let

$$\sigma_1 \ge \sigma_2 \ge \cdots \ge \sigma_0 \ge 0 \tag{21}$$

be the singular value decomposition of  $C_1$  where  $U_1$ ,  $V_1$  are square unitary matrices. (20) can be written equivalently

max tr 
$$V_1^t S_1 S_1^t V_1$$
 diag  $\left(\frac{\sigma_1^2}{1+\sigma_1^2}, \dots, \frac{\sigma_p^2}{1+\sigma_p^2}, 0, \dots, 0\right)$ 

subject to: V : unitary

$$0 \le \sigma_1 < a$$
,  $0 \le \sigma_2 \le a$ ,...,  $0 \le \sigma_0 \le a$ 

We obviously choose  $\sigma_1 = \sigma_2 = \cdots = \sigma_\rho = a$ . We also set

$$V_1 = [V_{11} \mid V_{12}]$$

$$V_{11} = [v_1 v_2 \dots v_p]$$
(23)

and we thus obtain

max tr 
$$V'_{11}S'_{1}S_{1}V_{11}$$
 (24)  
subject to:  $V'_{11}V_{11} = I$ 

To solve (24), we just choose  $v_1, v_2, \ldots, v_{\rho}$  to be the (orthonormalized) eigenvectors of  $S_1^i S_1$  corresponding to its  $\rho$  largest eigenvalues  $\lambda_1, \ldots, \lambda_{\rho}$ . Having found

 $V_{11}$  and  $\Sigma$  we can go back to (21) and construct C, using any arbitrary  $U_1$ .  $V_{12}$  can be taken to be any matrix which adjoined to  $V_{11}$  yields  $V_1$  unitary, but its choice does not matter as  $V_{12}'$  will be nullified when  $V_1'$  is multiplied by  $\Sigma$ . Notice that the optimal C is

$$C_1^* = \frac{a^2}{1+a^2} \ U V_1' \tag{25}$$

and the optimal cost

$$J^{*}(C_{1}^{*}) = -\frac{1}{2} \frac{a^{2}}{1+a^{2}} \operatorname{tr} V_{11}^{'} S_{1}^{'} S_{1}^{'} V_{11}$$
$$= -\frac{1}{2} \frac{a^{2}}{1+a^{2}} (\lambda_{1} + \lambda_{2} + \dots + \lambda_{\rho})$$
(26)

If there is no constraint on C, i.e.,  $a \rightarrow +\infty$  we obtain

$$C_1^* = VV_1'$$

$$J^*(C^*) = -\frac{1}{2}(\lambda_1 + \lambda_2 + \dots + \lambda_p)$$

in agreement with the results concerning the problem (15), (16) where  $C_1, S_1$  play now the role of C,S. Finally notice that the interesting feature of this problem is the separation between the magnitude

Finally notice that the interesting feature of this problem is the separation between the magnitude and rank constraints on C<sub>1</sub>.

# 4. Case 2: Two Decision Makers with Restricted Number of Measurements

In this section we consider the problem

minimize 
$$\hat{J}(C_1, C_2)$$
  
subject to: rank  $(C_1) \leq \rho_1$ 

$$rank (C_2) \le \rho_2 \tag{27}$$

The rank condition on  $C_i$  represents the inability of the decision maker i to acquire or process more than  $\rho_i$  measurements. If  $m_i \leq \rho_i$ , then the constraint rank  $(C_i) \leq \rho_i$  can be deleted, since the decision maker i does not really need more than  $m_i$  measurements. We can thus assume that  $\rho_i \leq n$ ,  $\rho_i \leq m_i-1$ . We can also substitute the inequality constraints in (35) with equality constraints, since more information does not hurt. For fixed  $C_1$ ,  $C_2$ , the optimal  $\gamma_1^*$ ,  $\gamma_2^*$  which minimize  $J(\gamma_1, \gamma_2)$  are given by (5)-(7). Multiplying (6) and (7) from the left by  $C_1'(C_1C')^{-1}$ ,  $C_2'(C_2C_2')^{-1}$  respectively, yields

$$\begin{split} & L_1 + RL_2C_2C_1'(C_1C_1')^{-1} + S_1C_1'(C_1C_1')^{-1} = 0 \\ & L_2 + R'L_1C_1C_2'(C_2C_2')^{-1} + S_2C_2'(C_2C_2') = 0 \end{split}$$

We can also assume, without loss of generality, that  $C_i C_i' = \text{unit}$ , since we can premultiply each  $y_i$  by some matrix. Let

$$C_1 C_2' = \Sigma (\rho_1 \times \rho_2)$$
 (28)

We thus have the system

$$L_1 + RL_2\Sigma' + S_1C_1' = 0 (29)$$

$$L_2 + R'L_1\Sigma + S_2C_2' = 0 (30)$$

A similar system can be obtained in the case of more than two decision makers. In the case of two decision makers, we can also assume without loss of generality that

$$\Sigma = \begin{bmatrix} \sigma_{1} & 0 & 0 & \cdots & 0 & \cdots & 0 \\ 0 & \sigma_{2} & 0 & \cdots & & & \vdots & \vdots \\ & & & \vdots & & \vdots & \vdots \\ 0 & & & 0 & \sigma_{\rho_{1}} & 0 & \cdots & 0 \end{bmatrix}$$

$$\rho_{1} \leq \rho_{2}$$

$$1 \geq \sigma_{1} \geq \cdots \geq \sigma_{\rho_{1}} \geq 0$$
(31)

(see [6]), so that the system (29), (30) can be easily solved explicitly for  $L_1$ ,  $L_2$ . (Unfortunately, such a simplification is not in general possible for the case of three or more decision makers.)

Let
$$L_{1} = [t_{1}, \dots, t_{\rho_{1}}]$$

$$L_{2} = [\overline{t}_{1} \dots \overline{t}_{\rho_{2}}]$$

$$C'_{1} = [v_{1} \dots v_{\rho_{1}}]$$
(32)

$$C_2' = [\overline{v}_1 \dots \overline{v}_{\rho_2}]$$
.

(29) and (30) can be written equivalently as

$$\begin{array}{lll}
\boldsymbol{L}_{i} + \boldsymbol{\sigma}_{i} \mathbf{R} \, \overline{\boldsymbol{L}}_{i} + \boldsymbol{S}_{1} \mathbf{v}_{i} &= 0 \\
\overline{\boldsymbol{L}}_{i} + \boldsymbol{\sigma}_{i} \mathbf{R}' \, \boldsymbol{L}_{i} + \boldsymbol{S}_{2} \overline{\mathbf{v}}_{i} &= 0 \\
\overline{\boldsymbol{L}}_{j} + \boldsymbol{S}_{2} \overline{\mathbf{v}}_{j} &= 0 , \quad j = \rho_{1} + 1, \dots, \rho_{2}
\end{array} \tag{33}$$

We can thus solve for  $L_i$ ,  $\overline{L}_i$ : substitute  $u_i = L_i y_i$  in the cost and find

$$\hat{\mathbf{J}}(C_{1},C_{2}) = -\frac{1}{2} \sum_{i=1}^{\rho_{1}} \begin{bmatrix} S_{1} & \mathbf{v}_{i} \\ S_{2} & \overline{\mathbf{v}}_{i} \end{bmatrix}' \begin{bmatrix} \mathbf{I} & \sigma_{i} \mathbf{R} \\ \sigma_{i} \mathbf{R}' & \mathbf{I} \end{bmatrix}^{-1} \begin{bmatrix} S_{1} & \mathbf{v}_{i} \\ S_{2} & \overline{\mathbf{v}}_{i} \end{bmatrix}$$

$$-\frac{1}{2} \sum_{j=\rho_{1}+1}^{\rho_{2}} \overline{\mathbf{v}}_{j}' S_{2}' S_{2} \overline{\mathbf{v}}_{j}$$
(34)

Thus, solving (27) is equivalent to solving the following problem:

subject to: 
$$\|\mathbf{v}_i\| = 1$$
,  $i = 1, ..., \rho_1$   
 $\|\overline{\mathbf{v}}_i\| = 1$ ,  $i = 1, ..., \rho_2$   
 $\mathbf{v}_i^t \mathbf{v}_j = \mathbf{v}_i^t \overline{\mathbf{v}}_j = \overline{\mathbf{v}}_j^t \overline{\mathbf{v}}_j = 0$ ,  $i \neq j$   
 $\mathbf{v}_i^t \overline{\mathbf{v}}_i = \sigma_i$ ,  $i = 1, 2, ..., \rho_1$ 

The general solution of (35) seems, at present, hard to come by. As an example, let us consider the case where  $\rho_1 = \rho_2$ ,  $R = \mu I$ ,  $S_1 = s_1 I$ ,  $S_2 = s_2 I$ . Then (43) assumes the form

$$\max_{i=1}^{\rho_1} \frac{s_1^2 + s_2^2 - 2\mu s_1 s_2 \sigma_i^2}{1 - \mu^2 \sigma_i^2}$$
(36)

subject to:  $0 \le \sigma_i \le 1$ ,  $i = 1, ..., \rho_1$ 

Since 2 each term in the summation (36), for  $0 \le \sigma_2 \le 1$ , is a piece of hyperbola, to maximize (36) we have to have:

$$\sigma_{i} = 1$$
 if  $\frac{s_{1}^{2} + s_{2}^{2} - 2\mu s_{1}s_{2}}{1 - \mu^{2}} \ge s_{1}^{2} + s_{2}^{2}$ 

Thus if  $\mu^2(s_1^2+s_2^2)-2\mu s_1s_2\geq 0$  we can choose  $C_1=C_2$  = any matrix with rank  $\rho_1$ .

If  $\mu^2(s_1^2+s_2^2)-2\mu s_1s_2<0$  and  $\rho_1+\rho_2=2\rho_1\leq n$  we can choose  $\sigma_i=0$ ,  $i=1,\ldots,\rho_1$  and thus choose  $C_1,C_2$  to be any two matrices of rank  $\rho_1=\rho_2$  such that  $C_1C_2=0$ . The only difficulty appears if  $\rho_1+\rho_2=2\rho_1>n$  and  $\rho^2(s_1^2+s_2^2)-2\rho s_1s_2<0$ , since we cannot now have  $C_1C_2=0$  and rank  $C_1=rank$   $C_2=\rho_1=\rho_2$ . A little reflection will persuade the reader that in this case we can choose  $C_1'=[v_1,\ldots,v_{\rho_1}]$   $C_2'=[\overline{v_1},\overline{v_2},\ldots,\overline{v_{\rho_1}}]$  is any orthonormal basis of  $R^n$  and  $\overline{v_{n-\rho_1}}+1,\ldots,\overline{v_{\rho_1}}$  have the same span as  $v_{n-\rho_1}+1,\ldots,v_{\rho_1}$ . As another example, consider the case where R,  $S_1$ ,  $S_2$  are  $2\times 2$  matrices, and  $\rho_1=\rho_2=1$ . By transforming  $v_1$  to  $V_1v_1$  where  $V_1,V_2$  are unitary matrices and  $v_1=v_2$  are unitary matrices and  $v_2=v_1$  diag  $(\mu_1,\mu_2)V_2$  is the singular value decomposition of  $v_1$ , with  $v_2=v_1$  we can consider without loss of generality that  $v_1=0$  are vectors on the plane we can set

$$\mathbf{v}_{1} = \begin{bmatrix} \cos \varphi \\ \sin \varphi \end{bmatrix}, \quad \mathbf{v}_{1} = \begin{bmatrix} \cos \theta \\ \sin \theta \end{bmatrix}, \quad \sigma_{1} = \cos (\varphi - \theta)$$

and solve an unconstrained problem with unknowns  $\varphi$ ,  $\theta$ . (We can actually restrict our attention to the square  $-\pi/2 \le \varphi \le \pi/2$ .) The objective function of this problem is the sum of two terms, each one of which is a quotient with nominator sums of powers of  $\cos \theta$ ,  $\sin \theta$  and denominator  $1-\mu_2^2 \cos^2(\varphi-\theta)$ .

Let us now consider a different approach for problem (27). Since for given  $C_i$ ,  $u_i$  will be linear in  $y_i$  and thus in x (recall also (19)-(21)) we set

$$u_1 = X_1 x \qquad u_2 = X_2 x \qquad (37)$$

where  $X_i$  is an  $m_i \times n$  matrix and consider the equivalent to (35)

$$\min \tilde{J}(X_1, X_2) = \operatorname{tr} \left[ \frac{1}{2} X_1' X_1 + \frac{1}{2} X_2' X_2 + X_1' R X_2 + X_1' S_1 + X_2' S_2 \right]$$
(38)

subject to: rank  $X_1 \le \rho_1$ 

$$rank X_2 \leq \rho_2$$

If  $X_2$  is fixed, the minimization in (38) with respect to  $X_1$  can be carried out by solving

min tr 
$$\left[\frac{1}{2}(X_1+RX_2+S_1)'(X_1+RX_2+S_1) - \frac{1}{2}(RX_2+S_1)'(RX_2+S_1) + X_2'S_2\right]$$
 (39)

subject to: rank  $X_1 \leq \rho_1$ 

To solve (39) for  $X_1$  we consider the singular value decomposition of  $RX_2+S_1$ :

$$RX_2 + S_1 = U_1' \Sigma_1 V_1$$
 (40)

where  $\text{U}_1,\ \text{V}_1$  are square unitary matrices, and  $\Sigma_1$  is of the form

$$\Sigma_{1} = \begin{bmatrix} \sigma_{1}^{1} & O \\ \sigma_{2}^{1} & O \\ & \sigma_{3}^{1} \\ O & \ddots \end{bmatrix}_{(m_{1} \times n)}, \ \sigma_{1}^{1} \ge \sigma_{2}^{1} > \dots > 0 \quad (41)$$

The X<sub>1</sub> that solves (48) is given by

$$\overline{X}_1 = U_1^{\prime} \overline{\Sigma}_1 V_1 \tag{42}$$

where

Similarly, for fixed  $X_1$ , to find  $X_2$  we consider the singular value decomposition

$$R'X_1 + S_2 = U_2'\Sigma_2V_2$$

$$\Sigma_{2} = \begin{bmatrix} \sigma_{1}^{2} & O \\ \sigma_{2}^{2} & O \\ O & \sigma_{3}^{2} & O \end{bmatrix} \qquad \sigma_{1}^{2} \ge \sigma_{2}^{2} \ge \sigma_{3}^{2} > \cdots > 0 \qquad (44)$$

$$(m_{2} \times n_{1})$$

and choose

$$\overline{X}_{2} = U_{2}^{1} \Sigma_{2} V, \quad \overline{\Sigma}_{2} = \begin{bmatrix} -\sigma_{1}^{2} & & & & \\ -\sigma_{2}^{2} & & & & \\ & -\sigma_{2}^{2} & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\$$

Thus, the problem is reduced to the following: Out of all pairs  $X_1$ ,  $X_2$  which satisfy (90)-(95) choose the one that results to the minimum value of the objective function in (38). If this pair is  $(X_1^*, X_2^*)$  we can choose  $C_1 = X_1^*$ ,  $C_2 = X_2^*$  and the  $L_1$ ,  $L_2$  of (6), (7) can be taken to be unit matrices.

This last way of tackling the problem, although is interesting, does not facilitate very much the solution of the problem, as long as we do not at present know how to find explicitly all the pairs  $X_1, X_2$  which satisfy (40)-(45). It nonetheless suggests an algorithm for generating pairs  $X_1, X_2$ 

which satisfy (40)-(45), namely the following: For fixed  $X_2 = X_2^0$  solve

$$\min \tilde{J}(X_1, X_2^0)$$

$$rank(X_1) \leq \rho_1$$

according to (40)-(43). Let  $X_1^0$  be the solution. Next for fixed  $X_1 = X_1^0$ , solve

$$\min \tilde{J}(X_1^0, X_2)$$

$$\operatorname{rank} X_2 \leq \rho_2$$

according to (44)-(45). Let  $X_2^1$  be the solution. Fix  $X_2 = X_2^1$  and generate  $X_1^1$  and so on. This algorithm is essentially the same with the one suggested in [1]. It obviously holds:

$$\tilde{\mathtt{J}}(x_1^{k+1}, \, x_2^{k+1}) \leq \tilde{\mathtt{J}}(x_1^{k+1}, \, x_2^k) \leq \tilde{\mathtt{J}}(x_1^{\,k}, \, x_2^k) \leq \cdots \leq \tilde{\mathtt{J}}(x_1^0, x_2^0)$$

and  $\tilde{J}$  is bounded from below by the best cost  $J^{**}$  corresponding to the case  $y_1=y_2=x$ . It is easy to verify that because  $\tilde{J}(X_1,X_2)$  is a quadratic and strictly convex function of  $X_1,X_2$ , the set of  $(X_1,X_2)$  which satisfy  $\tilde{J}(X_1,X_2) \leq \tilde{J}(X_1^0,X_2^0) = \text{constant}$  is a compact set, so that the sequence  $(X_1^k,X_2^k)$  has

necessarily at least one convergent subsequence. Thus the algorithm just described is guaranteed to provide in the limit a pair  $(X_1,X_2)$  which satisfies the relations (40)-(43). Unfortunately, this limit is not guaranteed to be the solution of problem (27) or equivalently of (38). [Notice that problem (27) is guaranteed to have a solution as  $J(C_1,C_2)$  is a continuous function of  $C_1,C_2$ , which  $C_1,C_2$  are assumed to be  $\rho_i \times n$  matrices, on which a magnitude constraint of the form  $\|C_i\| \leq a_i \ (a_1,a_2>0)$  can be imposed without loss of generality, since what matters is only the ranges of  $C_1',C_2'$ . Thus (27) can be considered as a problem of minimizing a continuous function subject to compact constraints and thus has a global solution.]

The following example demonstrates that this algorithm might fail to converge to the global solution of the problem.

### Example

Let 
$$n = m_1 = m_2 = 2$$
,  $\rho_1 = \rho_2 = 1$ 

$$R = \begin{bmatrix} \mu_1 & 0 \\ 0 & \mu_2 \end{bmatrix}, S_1 = \begin{bmatrix} s_1 & 0 \\ 0 & s_2 \end{bmatrix}, S_2 = \begin{bmatrix} \overline{s}_1 & 0 \\ 0 & \overline{s}_2 \end{bmatrix}$$

We will consider pairs  $X_1$ ,  $X_2$  which satisfy (40)-(43) and we will restrict our attention to  $X_1$ ,  $X_2$  diagonal, i.e.,

$$X_{1} = \begin{bmatrix} x_{1} & 0 \\ 0 & x_{2} \end{bmatrix}, \quad X_{2} = \begin{bmatrix} y_{1} & 0 \\ 0 & y_{2} \end{bmatrix}$$

(Notice that although nondiagonal  $X_1$ ,  $X_2$  which satisfy (40)-(43) might exist, if we choose  $X_2$  diagonal, and find  $X_1$  according to (40)-(43),  $X_1$  will also be diagonal.) It holds

$$RX_{2} + S_{1} = \begin{bmatrix} \mu_{1}y_{1} + s_{1} & 0 \\ 0 & \mu_{2}y_{2} + s_{2} \end{bmatrix}$$

$$R'X_{1} + S_{2} = \begin{bmatrix} \mu_{1}x_{1} + \overline{s}_{1} & 0 \\ 0 & \mu_{2}x_{2} + \overline{s}_{2} \end{bmatrix}$$

We have the following four cases:

(a) 
$$x_1 = (\mu_1 \overline{s}_1 - s_1)/(1 - \mu_1^2)$$
  $x_2 = 0$   
 $y_1 = (\mu_1 s_1 - \overline{s}_1)/(1 - \mu_1^2)$   $y_2 = 0$   
if  $\begin{cases} |x_1| \ge |s_2| \\ |y_1| \ge |\overline{s}_2| \end{cases}$ 

(a) 
$$x_1 = 0$$
  $x_2 = (\mu_2 \overline{s}_2 - s_2)/(1 - \mu_2^2)$ 

$$y_1 = 0$$
  $y_2 = (\mu_2 \overline{s}_2 - \overline{s}_2)/(1 - \mu_2^2)$ 
if  $\begin{cases} |s_1| \le |x_2| \\ |\overline{s}_1| \le |y_2| \end{cases}$ 

$$\begin{aligned} & (\gamma) & & \mathbf{x}_1 = -\mathbf{s}_1 & & \mathbf{x}_2 = 0 \\ & & \mathbf{y}_1 = 0 & & \mathbf{y}_2 = -\overline{\mathbf{s}}_2 \\ & & & \left\{ \begin{vmatrix} \mathbf{s}_1 \end{vmatrix} \geq |\mu_2 \overline{\mathbf{s}}_2 - \mathbf{s}_2 \rangle \\ |\mu_1 \mathbf{s}_1 - \overline{\mathbf{s}}_1| \leq |\overline{\mathbf{s}}_2| \end{vmatrix} \right. \end{aligned}$$

(6) 
$$x_1 = 0$$
  $x_2 = -s_2$ 

$$y_1 = -\overline{s}_1$$
  $y_2 = 0$ 
if  $\begin{cases} |\mu_1 \overline{s}_1 - s_1| \le |s_2| \\ |\overline{s}_1| \ge |\mu_2 s_2 - \overline{s}_2| \end{cases}$ 

It is easy to see that there are choices for  $\mu_1$ ,  $\mu_2$ ,  $s_1$ ,  $s_2$ ,  $s_1$ ,  $s_2$  so that more than one of the cases  $\alpha$ ,  $\beta$ ,  $\gamma$ ,  $\delta$  are acceptable. For example, if  $\mu_1 = \mu_2 = 1/2$ ,  $s_1 = -7/2$ ,  $s_2 = 4$ ,  $s_1 = 0$ ,  $s_2 = 2$ , then both cases  $\gamma$  and  $\delta$  are acceptable.

Thus the algorithm described might fail to reach the global optimum of (27). Actually, to solve the above mentioned example, within the class of diagonal  $X_1, X_2$ , we have to check which out of the four cases  $\alpha, \beta, \gamma, \delta$  are acceptable and if more than one is, to calculate the values of  $\hat{J}$  at each one of them and choose the one which results to the smallest.

There is a third, slightly different, approach that one could follow for solving the problem. Problem (27) can be written equivalently as follows:

min tr 
$$\left[\frac{1}{2}C_{1}^{\prime}L_{1}^{\prime}L_{1}C_{1}+\frac{1}{2}C_{2}^{\prime}L_{2}^{\prime}L_{2}C_{2}+C_{1}^{\prime}L_{1}^{\prime}RL_{2}^{\prime}C_{2}+C_{1}^{\prime}L_{1}^{\prime}RL_{2}^{\prime}C_{2}+C_{1}^{\prime}RL_{2}^{\prime}C_{2}+C_{1}^{\prime}RL_{2}^{\prime}C_{2}+C_{1}^{\prime}RL_{2}^{\prime}C_{2}+C_{1}^{\prime}RL_{2}^{\prime}C_{2}+C_{1}^{\prime}RL_{2}^{\prime}C_{2}+C_{1}^{\prime}RL_{2}^{\prime}C_{2}+C_{1}^{\prime}RL_{2}^{\prime}C_{2}+C_{1}^{\prime}RL_{2}^{\prime}C_{2}+C_{1}^{\prime}RL_{2}^{\prime}C_{2}+C_{1}^{\prime}RL_{2}^{\prime}C_{2}+C_{1}^{\prime}RL_{2}^{\prime}C_{2}+$$

subject to: 
$$L_1 + RL_2C_2C_1' + S_1C_1' = 0$$
 (46-1)

$$L_2 + R'L_1 C'_1 C'_2 + S_2 C'_2 = 0$$
 (46-2)

$$C_1 C_1' = I(\rho_1 \times \rho_1)$$
 (46-3)

$$C_2C_2' = I(\rho_1 \times \rho_2)$$
 (46-4)

Notice that  $L_1$ ,  $L_2$ ,  $C_1$ ,  $C_2$  are considered as unknowns. It can be easily verified that a Lagrange multiplier by which we can append the constraints exists, by the following argument: (46-1) and (46-2) are always uniquely solvable for  $L_1$ ,  $L_2$ , given  $C_1$  and  $C_2$  as in (46-3) and (46-4); one can also verify that the full rank condition with respect to the unknowns  $C_1$ ,  $C_2$  is always fulfilled by the gradient of (46-3) and (46-1); thus the full rank condition is satisfied by the constraints and consequently a Lagrange multiplier exists. We now append the constraints to the objective, take the gradient to be zero and after some calculations we end up with the following necessary conditions that have to be satisfied by  $L_1$ ,  $L_2$ ,  $C_1$ ,  $C_2$ :

$$L_1 + RL_2C_2C_1' + S_1C_1' = 0 (47-1)$$

$$L_2 + R'L_1C_1C_2 + S_2C_2 = 0 (47-2)$$

$$L'_1(L_1C_1+RL_2C_2+S_1)=0$$
 (47-3)

$$L'_{2}(L'_{2}C_{2} + R'L_{1}C_{1} + S_{2}) = 0$$
 (47-4)

$$C_1 C_1' = I$$
 (47-5)

$$C_2C_2' = I$$
 . (47-6)

(47-3) can be multiplied from the left by  $C_1'$  to yield equivalently

$$C'_1L'_1(L_1C_1+RL_2C_2+S_1)=0$$
 (48)

Notice that multiplying (48) from the left by  $C_1$  yields (47-3) because of (47-5). It is obvious now that by setting  $X_i = L_i C_i$ , (47) can be written equivalently

$$X_1 + RX_2C_1'C_1 + S_1C_1'C_1 = 0$$
 (49-1)

$$X_2 + R'X_1C_2C_2 + S_2C_2C_2 = 0$$
 (49-2)

$$X'_1(X_1+RX_2+S_1)=0$$
 (49-3)

$$X_2'(X_2+R'X_1+S_2)=0$$
 (49-4)

$$C_1 C_1' = I(\rho_1 \times \rho_1)$$
 (49-5)

$$C_2 C_2' = I(\rho_2 \times \rho_2)$$
 (49-6)

A little reflection will persuade the reader that (49-3) could have been directly derived from (40)-(43) and similarly (49-4) from (44)-(45). One could in principle solve (47-1), (47-2) explicitly for  $L_1$ ,  $L_2$ , plug their values in (47-3), (47-4) and have a system of equations that should be satisfied by  $v_1, \dots, v_{\rho_1}, \overline{v_1}, \dots, \overline{v_{\rho_2}}$ . One could do the same thing with (49-1)-(49-4). (49-1)-(49-2) are easy to solve explicitly for  $X_1, X_2$  under the assumption that the R matrix is square and  $R = \mu I$  for some  $\mu$ :  $|\mu| < 1$ , in which case (49-1)-(49-6) can be simplified (after some calculations) to the equivalent:

$$[\mu C_1 C_2^{'} C_2 S_2^{'} S_2 C_2^{'} - C_1 S_1^{'} S_2 C_2^{'}][I - C_2 C_1^{'} C_1 C_2^{'}] = 0 \quad (50)$$

$$[\mu C_2 C_1' C_1 S_1' S_1 C_1' - C_2 S_2' S_1 C_1'][I - C_1 C_2' C_2 C_1'] = 0. (51)$$

(50) and (51) characterize the optimal C1, C2.

Example

Let x be two-dimensional,  $\rho_1 = \rho_2 = 1$  and  $R = \mu I$ . The unknowns  $C_1$ ,  $C_2$  can be taken to be

$$C_1 = [\cos \varphi, \sin \varphi], \quad C_2 = [\cos \theta, \sin \theta].$$

It holds

$$C_1C_2' = \cos(\varphi - \theta)$$

If

$$\cos (\varphi - \theta) \neq \pm 1$$
,

(50) and (51) yield a system of two equations with unknowns,  $\phi$  and  $\theta$ , which can be simplified to the form

$$\mu \cos (\varphi - \theta) \|S_2 \begin{bmatrix} \cos \theta \\ \sin \theta \end{bmatrix}\|^2$$

= 
$$[\cos \varphi, \sin \varphi] S_1' S_2 \begin{bmatrix} \cos \theta \\ \sin \theta \end{bmatrix}$$
 (52)

$$\|S_{2}\begin{bmatrix}\cos\theta\\\sin\theta\end{bmatrix}\|^{2} = \|S_{1}\begin{bmatrix}\cos\varphi\\\sin\phi\end{bmatrix}\|^{2}$$
 (53)

The geometrical meaning of these two conditions is that the vectors

$$S_{1}\begin{bmatrix}\cos\phi\\\sin\phi\end{bmatrix}$$
,  $S_{2}\begin{bmatrix}\cos\theta\\\sin\theta\end{bmatrix}$ 

have equal lengths and their angle w satisfies

$$\cos w = \mu \cos (\varphi - \theta)$$
.

The case cos  $(\phi - \theta) = \pm 1$ , i.e.,  $\phi = \theta + integer$  multiple of  $\pi$  --or without loss of generality:  $\phi - \theta$  -- can be easily examined separately.

5. Case iii. Two Decision Makers with Restricted
Number of Measurements and No
Common Information

In this section we consider the problem

$$\min \hat{J}(C_1, C_2)$$

subject to: rank 
$$(C_1) \le \rho_1$$
 (54)  
rank  $(C_2) \le \rho_2$   
 $C_1 C_2' = 0$ 

The meaning of the additional condition  $C_1C_2'=0$  (compare with (27) is that there is no common information between the two decision makers. In this case the term  $u_1'Ru_2$  is of no importance, since its expectation is obviously zero. Use of (6)-(7) yields that (54) is equivalent to solving

min tr 
$$[-\frac{1}{2}S_1'S_1P_1 - \frac{1}{2}S_2'S_2P_2]$$

subject to:

rank 
$$(P_1) \le \rho_1$$
, rank  $(P_2) \le \rho_2$   
 $P_1P_2 = 0$ 

Recall that:

$$P_{i} = C_{i}(C_{i}C_{i}^{i})^{+}C_{i}$$
,  $i = 1, 2$ .

(55) is the obvious generalization of (16).

Two main cases can now be considered. The first case is  $\rho_1 + \rho_2 \ge n$  and the second one  $\rho_1 + \rho_n < n$ . The first case is quite easy to solve as the following argument shows: we can take rank  $(C_1)$  + rank  $(C_2)$  = n in which case  $P_2$  = I -  $P_1$ , i.e., we have only one unknown. In particular, let

rank 
$$(P_1) = n-1$$
, rank  $(P_2) = n - \rho_1 + 1$ ,

where & is an integer satisfying:

$$0 \le \ell \le \rho_1 + \rho_2 - n .$$

For fixed £ (55) assumes the form

max tr 
$$[(S_1^tS_1 - S_2^tS_2)P_1]$$
  
subject to: rank  $P_1 = n-t \le P_1$ 

We obviously choose  $\ell = (n-\rho_1)$  and we solve

max tr 
$$[(S_1'S_1 - S_2'S_2)P_1]$$
  
subject to: rank  $(P_1) = \rho_1$ 

which falls within the class of problems solved in Section 2 and can be solved, by taking  $P_1$  to be the projection on the space spanned by the  $\rho_1$  eigenvectors of  $S_1^{\mathsf{I}}S_1^{\mathsf{I}}-S_2^{\mathsf{I}}S_2^{\mathsf{I}}$  corresponding to the  $\rho_1$  largest eigenvalues. So, we only need to concentrate on the case  $\rho_1^{\mathsf{I}}+\rho_2^{\mathsf{I}}<\infty$ .  $P_i$  can be taken to be

$$P_{i} = U_{i}U_{i}^{!}$$
,  $i = 1, 2$ 

where  $U_i$  is an  $n \times \rho_i$  matrix with  $U_i^! U_i^! = \text{unit } \rho_i \times \rho_i$  and  $U_1^! U_2^! = \text{zero matrix}$ . Problem (55) assumes the form

max 
$$U_{1}^{\prime}S_{1}^{\prime}S_{1}U_{1}^{\prime} + U_{2}^{\prime}S_{2}^{\prime}S_{2}U_{2}^{\prime}$$
  
subject to:  $U_{1}^{\prime}U_{1} = I(\rho_{1} \times \rho_{1})$  (56)  
 $U_{2}^{\prime}U_{2} = I(\rho_{2} \times \rho_{2})$   
 $U_{1}^{\prime}U_{2} = O(\rho_{1} \times \rho_{2})$ 

Because of the compactness of the constraint set and the continuity of the objective function, problem (54) has obviously a solution. Appending the constraints with Lagrange multipliers—which exist because the constraints satisfy the full rank condition as can be easily verified—we obtain the following necessary conditions for (54):

$$S_1 S_1 U_1 = U_1 \Lambda_1 + U_2 \Lambda$$
 (57)

$$S_{2}^{\prime}S_{2}U_{2} = U_{2}\Lambda_{2} + U_{1}\Lambda^{\prime}$$
 (58)

$$U_1U_1 = I$$

$$U_2'U_2 = I (59)$$

$$U_1^{\dagger}U_2 = 0$$

where  $\Lambda_1$  and  $\Lambda_2$  are symmetric matrices. Out of all the  $U_1, U_2, \Lambda_1, \Lambda_2, \Lambda$  that satisfy (57)-(59) we want the one that yields the maximum falue for

$$tr(\Lambda_1+\Lambda_2)$$
.

(57)-(59) is a "generalized" eigenvalue type of problem. As an example, let us consider the case where  $\rho_1 = \rho_2 = 1$ , and  $S_1'S_1, S_2'S_2$  are 3x3 matrices.

Then (57)-(59) reduce to finding vectors v<sub>1</sub>, v<sub>2</sub> such that

$$S_{1}^{i}S_{1}v_{1} = \lambda_{1}v_{1} + \lambda v_{2}$$

$$S_{2}^{i}S_{2}v_{2} = \lambda_{2}v_{2} + \lambda v_{1}$$

$$v_{1}^{i}v_{1} = v_{2}^{i}v_{2} = 1$$

$$v_{1}^{i}v_{2} = 0$$
(60)

For the case where  $S_1'S_1$ ,  $S_2'S_2$  are  $3\times3$  matrices, and  $\rho_1 = \rho_2 = 1$ , we can solve the problem in the following way. If we knew the vector  $\mathbf{v}_3$  which spans the space perpendicular to  $\mathbf{v}_1$  and  $\mathbf{v}_2'$  (i. e.,  $\mathbf{v}_2'$ ,  $\mathbf{v}_1 = \mathbf{v}_3'\mathbf{v}_2 = 0$ ,  $\mathbf{v}_3'\mathbf{v}_3 = 1$ ) then we could solve the problem

$$\max \operatorname{tr} (I - P_3)S_1S_1'(I - P_3)P_1 + (I - P_3)S_2S_2'(I - P_3)P_2$$

$$P_1 = v_1v_1', \quad P_2 = v_2v_2'$$

$$P_1P_2 = 0$$
(61)

where the unknowns are  $v_1$ ,  $v_2$  and  $P_3 = v_3 v_3^t$  is known. This problem can be solved since rank  $P_1$ + rank  $P_2$  = 2 which is equal to the dimension of the space where we are working, i.e., the space where I -  $P_3$  project; (recall case with  $\rho_1 + \rho_2 = n$ ). Thus it suffices to be able to find  $v_3$ . It holds

$$tr [S'_1S_1P_1 + S'_2S_2P_2] =$$

$$= tr[(I-P_3)(S'_1S_1-S'_2S_2)(I-P_3)P_1 + S'_2S_2(I-P_3)]$$

Thus, we can consider equivalently to (61) the problem:

$$\max_{P_3} \left\{ \text{tr} \left( S_2' S_2 (I - P_3) \right) + \text{maximum eigenvalue} \right. \\ \left. \left[ (I - P_3) (S_1' S_1 - S_2' S_2) (I - P_3) \right] \right\}$$
 (62)

The maximum eigenvalue of  $(I-P_3)(S_1'S_1-S_2'S_2)(I-P_3) = \overline{\lambda}(v_3)$  can be found explicitly, for we are working with 3x3 matrices. Without loss of generality let

$$S_{1}^{\prime}S_{1}-S_{2}^{\prime}S_{2} = \begin{bmatrix} a_{1} & 0 & 0 \\ 0 & a_{2} & 0 \\ 0 & 0 & a_{3} \end{bmatrix}, a_{1} \ge a_{2} \ge a_{3} = 0 \quad (63)$$

$$\mathbf{v}_3 = \begin{bmatrix} \mathbf{n}_1 \\ \mathbf{n}_2 \\ \mathbf{n}_3 \end{bmatrix}, \quad \mathbf{n}_1^2 + \mathbf{n}_2^2 + \mathbf{n}_3^2 = 1.$$

To find  $\overline{\lambda}(v_3)$  we have to solve:

$$\det \left[ (\mathbf{I} - \mathbf{v}_3 \mathbf{v}_3') \begin{bmatrix} \mathbf{a}_1 & 0 & 0 \\ 0 & \mathbf{a}_2 & 0 \\ 0 & 0 & 0 \end{bmatrix} (\mathbf{I} - \mathbf{v}_3 \mathbf{v}_3') - \lambda \mathbf{I} \right] = 0$$

$$\lambda \left\{ \left[ a_{1}(1-n_{1}^{2})-\lambda \right] \left[ a_{2}(1-n_{2}^{2})-\lambda \right] - a_{1}a_{2}n_{1}^{2}n_{2}^{2} \right\} = 0$$

hue

$$\overline{\lambda}(v_3) = \frac{a_1(1-v_1^2) + a_2(1-v_2^2) + \sqrt{\left[a_1(1-v_1^2) + a_2(1-v_2^2)\right]^2 - 4a_1a_2v_3^2}}{2}$$

and (86) can be written equivalently as

$$\max_{\substack{n_1, n_2, n_3 \\ s. \text{ to: } n_1^2 + n_2^2 + n_3^2 = 1}} \left\{ -2 \begin{bmatrix} n_1 \\ n_2 \\ n_3 \end{bmatrix} : S_2^t S_2 \begin{bmatrix} n_1 \\ n_2 \\ n_3 \end{bmatrix} + a_1 (1-n_1^2) + a_2 (1-n_2^2) + a_1 (1-n_2^2) + a_2 (1-n_2^2) + a_2$$

If  $a_1 = a_2$ , then (64) assumes the form

min 
$$v_3'S_2'S_2v_3$$

and  $v_3$  will be chosen as the eigenvectors corresponding to the smallest eigenvalue of  $S_2^tS_2$ . If  $a_2=0$ , then (64) assumes the form

$$\min v_3'S_1'S_1v_3$$

and  $v_3$  will be chosen as the eigenvector corresponding to the smallest eigenvalue of  $S_1'S_1$ . If  $a_1>a_2>a_3$ , there is no obvious immediate simplification of (64), which has to be solved as a classical nonlinear programming problem. The only obvious conclusion is that if two of the eigenvalues of  $S_1'S_1-S_2'S_2$  coincide, then it is very easy to solve (64). If this is not the case, we can still exploit this previous conclusion in order to give upper and lower bounds for the optimal value. Let  $a_1>a_2>a_3$ . Then

$$S_{1}^{i}S_{1} = S_{2}^{i}S_{2} + \begin{bmatrix} a_{1} \\ a_{1} \\ a_{2} \\ a_{2} \end{bmatrix} = M_{2}$$

$$S_{2}^{i}S_{2} + \begin{bmatrix} a_{1} \\ a_{2} \\ a_{3} \end{bmatrix} = M_{2}$$

$$S_{1}^{i}S_{1} = S_{2}^{i}S_{2} + \begin{bmatrix} a_{1} \\ a_{2} \\ a_{3} \end{bmatrix} = N_{1},$$

$$S_{2}^{i}S_{2} + \begin{bmatrix} a_{1} \\ a_{2} \\ a_{3} \end{bmatrix} = N_{2}$$

$$S_{2}^{i}S_{2} + \begin{bmatrix} a_{1} \\ a_{3} \\ a_{3} \end{bmatrix} = N_{2}$$

and

$$\max \operatorname{tr}(S_{1}'S_{1}P_{1}+S_{2}'S_{2}P_{2}) \leq \max \operatorname{tr}(M_{1}P_{1}+S_{2}'S_{2}P_{2}) ,$$

$$i = 1, 2$$

$$\max \operatorname{tr}(S_{1}'S_{1}P_{1}+S_{2}'S_{2}P_{2}) \geq \max \operatorname{tr}(N_{1}P_{1}+S_{2}'S_{2}P_{2}) ,$$

$$i = 1, 2$$

The maximization problems on the right hand side of (65) can be easily solved to provide upper and lower bounds.

For the case where  $S_i^lS_i^l$  has dimension larger than 3 and  $P_i^l$  larger than 1, we can easily generalize several of the results presented and we can similarly create upper and lower bounds for the optimal values of the objective function.

Before completing this section let us consider another class of problems which can be reduced to those considered in this section. Consider the problem

min 
$$\hat{J}(C_1, C_2)$$
  
subject to: rank  $(C_1) \le \rho_1$   
rank  $(C_2) \le \rho_2$  (66)  
range  $(C_1) \subseteq \text{range}(C_2)$ 

In this case

$$P_1P_2 = P_2P_1 = P_1$$

 $(P_i = C_i (C_i C_i^i)^{\dagger} C_i^i)$  and (6), (7) can be explicitly solved for  $u_1, u_2$ , so that  $\hat{J}(C_1, C_2)$  can be explicitly calculated and is found to be of the form

$$\hat{J}(C_1, C_2) = -tr(A_1P_1 + A_2P_2)$$

where  $A_1$ ,  $A_2$  are known symmetric matrices which depend on R,  $S_1$ ,  $S_2$ . Let

$$\hat{P}_2 = P_2 - P_1 .$$

Ĵ assumes the form

$$\hat{J} = -tr[(A_1 + A_2)P_1 + A_2\hat{P}_2]$$

where  $P_1\hat{P}_2 = 0$  which shows that if we consider  $P_1, \hat{P}_2$  as unknowns we have reduced (66) to the form (55). It should be pointed out that the problem (66) is important in its own, since it is exactly the problem that has to be solved in a two stage dynamic linear quadratic gaussian problem, with no measurement noise, which is characterized by nested information.

# 6. <u>Case iv.</u> <u>Three Decision Makers with Independent Measurements</u>

Here we consider the case where the cost is given by

$$J = \frac{1}{2}(u_1^2 + u_2^2 + u_3^2) + u_1S_1'x + u_2S_2'x + u_3S_3'x$$

 $\mathbf{u}_1, \mathbf{u}_2, \mathbf{u}_3$  are scalar valued,  $\mathbf{S}_1, \mathbf{S}_2, \mathbf{S}_3$  are fixed fectors in  $\mathbf{R}^3$ , and x is a gaussian random vector in  $\mathbf{R}^3$  with zero mean and unit variance. The measurement available to  $\mathbf{u}_i$  is  $\mathbf{y}_i$  where

$$\mathbf{y_i} = \mathbf{v_i'x}$$
 ,  $\mathbf{v_i'v_j} = \delta_{ij}$  ,  $i, j = 1, 2, 3$  .

Notice that each decision maker's information is independent of the information of the others. The problem of optimal choice of information reduces to solving the following problem

$$\max v_{1}' S_{1} S_{1}' v_{1} + v_{2}' S_{2} S_{2}' v_{2} + v_{3}' S_{3} S_{3}' v_{3} =$$

$$= (S_{1}' v_{1})^{2} + (S_{2}' v_{2})^{2} + (S_{3}' v_{3})^{2}$$
(67)

subject to:  $v_i^! v_i = \delta_{ij}$ 

The geometrical interpretation of (67) is the following: Given three vectors,  $S_1$ ,  $S_2$ ,  $S_3$ , find the orthogonal parallelepipedon with maximal diagonal, whose one corner is at the origin and the opposite corner's three adjacent sides pass from the end points of the vectors  $S_1$ ,  $S_2$ ,  $S_3$ . The corresponding problem on the two dimensional plane is the following: Given two vectors  $S_1$ ,  $S_2$ , find the orthogonal parallelogrammon with maximal diagonal, whose one corner is at the origin and the opposite corner's two adjacent sides pass from the end points

of the vectors  $S_1$ ,  $S_2$ . The two dimensional case is easy to solve: If  $A_i$  is the end point of  $S_i$ , let M be the middle of  $A_1$ ,  $A_2$ . Consider the circle with diameter  $A_1A_2$  and let  $N_1$ ,  $N_2$  be the points where the line defined by OM meets the circle. Let  $OB_i$  be the perpendicular to  $N_1A_i$  from O. The parallelogrammon O  $B_1N_1B_2$  O is the one with maximal diagonal ( $N_2$  corresponds to the minimum). That the two dimensional case is easy to solve is not surprising, because it actually corresponds to the case where rank  $P_1$ + rank  $P_2$  = n,  $P_1P_2$  = 0 of Section 3. The three dimensional case that we are interested in corresponds to rank  $P_1$ + rank  $P_2$  = 1+1 < 3 = n.

Using the Lagrange multiplier rule for (67), yields the following necessary condition:

$$S_{1}S_{1}^{1}v_{1} = \lambda_{11}v_{1} + \lambda_{12}v_{2} + \lambda_{13}v_{3} \qquad \lambda_{12} = \lambda_{21}$$

$$S_{2}S_{2}^{1}v_{2} = \lambda_{21}v_{1} + \lambda_{22}v_{2} + \lambda_{23}v_{3} \qquad \lambda_{31} = \lambda_{13}$$

$$S_{3}S_{3}^{1}v_{3} = \lambda_{31}v_{1} + \lambda_{32}v_{2} + \lambda_{33}v_{3} \qquad \lambda_{23} = \lambda_{32}$$
(68)

(It is easy to justify the existence of a Lagrangle multiplier for this problem.) Out of all  $v_i$ ,  $\lambda_{ij}$  which satisfy (68) we are interested in the one which maximizes  $\lambda_{11}^{+}\lambda_{22}^{+}\lambda_{33}^{-}$ .

If  $S_1'v_1 = S_2'v_2 = 0$ , then (68) yields  $\lambda_{ij} = 0$  and thus  $\lambda_{11} + \lambda_{22} + \lambda_{33} = 0$ ; if at least one of the  $S_i$ 's is different than zero, then this contradicts the fact that the maximum in (67) has to be strictly positive. If  $S_1'v_1 = S_2'v_2 = 0$ , and  $S_3'v_3 \neq 0$ , then (68) yields  $\lambda_{11} = \lambda_{12} = \lambda_{13} = \lambda_{21} = \lambda_{22} = \lambda_{23} = \lambda_{31} = \lambda_{32} - 0$  and  $S_3'S_3'v_3 = \lambda_{33}v_3$  and thus  $v_3 = S_3/\|S_3\|$ . It can now be easily checked whether  $v_1, v_2$  exist with  $v_1'S_3 = v_2'S_3 = v_1'S_1 = v_2'S_2 = 0$ . If  $S_3'v_3 = 0$  and  $S_2'v_2 \neq 0$ ,  $S_1'v_1 \neq 0$  then (68) yields

$$S_1 S_1' v_1 = \lambda_{11} v_1 + \lambda_{12} v_2$$
  
 $S_2 S_2' v_2 = \lambda_{12} v_1 + \lambda_{22} v_2$ 

which actually means that we have to solve the two-dimensional analog discussed above. Similarly we can examine all the cases where at least one of the  $S_1^i v_1^i$ 's is zero. Thus we can concentrate on the case where  $S_1^i v_1^i$ ,  $S_2^i v_2^i$ ,  $S_3^i v_3^i \neq 0$ , in which case the conditions (68) can be written as

$$S = (S_1: S_2: S_3] = [v_1 v_2 v_3] \begin{bmatrix} \mu_1 & 0 & 0 \\ 0 & \mu_2 & 0 \\ 0 & 0 & \mu_3 \end{bmatrix} \begin{bmatrix} 1 & a & b \\ \alpha & 1 & c \\ b & c & 1 \end{bmatrix}$$
$$= UMR$$
 (69)

where  $\mu_i^2 = \lambda_{ii}$ . If  $S_1, S_2, S_3$  are linearly independent we can replace (69) equivalently by:

$$R(S'S)^{-1}R = diagonal$$

which yields a system of three equations with three unknowns. For each solution  $(\alpha, \beta, \gamma)$  of this system we can find the diagonal elements of  $R(S^iS)^{-1}R$  which are actually equal to  $1/(\mu_i)^2 = 1/\lambda_{ii}$  and pick the solution  $(\alpha, \beta, \gamma)$  that yields the maximum for  $\lambda_{11}^{+} + \lambda_{22}^{+} + \lambda_{33}^{+}$ . Having thus found  $\mu_1, \mu_2, \mu_3, a, b, c$ ,

we find  $U = SR^{-1}M^{-1}$ . The only difficulty in the above procedure lies in solving the system

$$\begin{bmatrix} 1 & a & b \end{bmatrix} (S'S)^{-1} \begin{bmatrix} a \\ 1 \\ c \end{bmatrix} = 0$$

$$\begin{bmatrix} 1 & a & b \end{bmatrix} (S'S)^{-1} \begin{bmatrix} b \\ c \\ 1 \end{bmatrix} = 0$$

$$\begin{bmatrix} a & 1 & c \end{bmatrix} (S'S)^{-1} \begin{bmatrix} b \\ c \\ 1 \end{bmatrix} = 0$$

for a, b, c.

#### 7. Conclusions

Our main objective in this paper was to formulate some problems related to the optimal choice of information in a team problem., Some partial results were also presented, which suggest several possible ways of handling these problems. As it turned out, several "generalized type eigenvalue" problems have to be solved and their geometric meaning to be connected in a simple and intuitive way with the matrices involved in describing the cost and information. We believe that the importance of the topic asks for further investigation of these issues.

### 8. References

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